

# Cycles in 3-connected Planar Graphs

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**PRE-CONFERENCE**

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# Introduction

- ▶ Let  $H \subset G$ ; then  $G/H$  denotes the graph with  $V(G/H) = V(G-H) \cup h$  and  $E(G/H) = E(G-H) \cup [hy: y \in V(G-H) \text{ and } yy' \in V(H)]$ . We say that  $G/H$  is obtained from  $G$  by contracting  $H$  to the vertex  $h$ .
- ▶ A block of a graph  $G$  is a maximal 2-connected subgraph of  $G$ .
- ▶ A cycle is a facial cycle if it bounds a face of the graph.
- ▶ Open, Closed Disc
- ▶ Let  $H$  be a subgraph of a graph  $G$ : An  $H$ -bridge of  $G$  is a subgraph of  $G$  which either (1) is induced by an edge of  $E(G) - E(H)$  with both incident vertices in  $H$  or (2) is induced by the edges in a component of  $G$  from  $H$  and the edges of  $G$  from  $H$  to that component. For any  $H$ -bridge  $B$  of  $G$ ; the attachments of  $B$  (on  $H$ ) are the vertices in  $V(B) \cap V(H)$ .

## Introduction-2

- ◀ A circuit graph is a pair  $(G,C)$ , where  $G$  is a 2-connected plane graph and  $C$  is a facial cycle of  $G$ ; such that, for any 2-cut  $S$  of  $G$ ; every component of  $G-S$  contains a vertex of  $C$ .
- ◀ An annulus graph is a triple  $(G,C_1,C_2)$  where  $G$  is a 2-connected plane graph and  $C_1$  and  $C_2$  are facial cycles of  $G$ ; such that, for any 2-cut  $S$  of  $G$ ; every component of  $G - S$  contains a vertex of  $C_1 \cup C_2$ .

# Introduction-Example

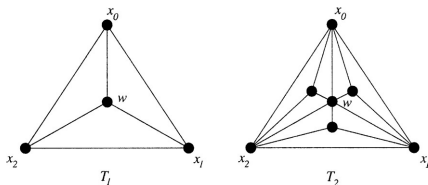


FIG. 1.

- ▶
- ▶ First, we define a sequence of 3-connected plane graphs  $T_k$  as follows. Let  $T_1$  be a plane graph isomorphic to  $K_4$ . Further, let  $V(T_1) = [w, x_0, x_1, x_2]$  and let  $x_0x_1x_2x_0$  be the outer cycle of  $T_1$ . Suppose that  $T_k$  is defined for some  $k \geq 1$ . Let  $T_{k+1}$  be the graph obtained from  $T_k$  as follows: In each inner face of  $T_k$ ; add a new vertex and join the new vertex to the vertices of  $T_k$  incident with that face. Graphs  $T_1$  and  $T_2$  are shown in Fig.1.

# Introduction-Example

- ◀ The above construction, for any  $k \geq 1$ ,  $T_k$  is a 3-connected plane graph with outer cycle  $x_0x_1x_2x_0$  with  $\text{circ}(T_k) < \frac{7}{2}n^{\log_3 2}$
- ◀ Let  $\alpha_k$  be the length of the longest  $x_1$ - $x_2$  path in  $T_k$  and  $\beta_k$  be the length of a longest  $x_1$ - $x_2$  path in  $T_k - x_0$ . By the construction of  $T_k$ , for  $i, j \in [0, 1, 2]$  and  $i \neq j$ , the length of a longest  $x_i - x_j$  path in  $T_k$  is  $\alpha_k$  and the length of a longest  $x_i - x_j$  path in  $T_k - ([x_0, x_1, x_2] - [x_i, x_j])$  is  $\beta_k$ .
- ◀ For  $i \in [0, 1, 2]$ , let  $D^i$  denote the open disc in the plane bounded by the triangle in  $T_{k+1}$  induced by  $[w, x_0, x_1, x_2] - [x_i]$ . Let  $V^i$  denote the set of vertices in  $T_{k+1}$  contained in  $D^i$ , and let  $T^i$  be the plane subgraph of  $T_{k+1}$  induced by  $V^i \cup ([w, x_0, x_1, x_2] - [x_i])$ .
- ◀ **Proposition** :  $\alpha_k = 3 * 2^{k-1}$  and  $\beta_k = 2^k$

# Introduction-Important Theorems

- ▶ Let  $(G, C)$  be a circuit graph, and let  $x, y \in V(C)$ . We say that  $(G, xCy)$  is a strong circuit graph if, for any 2-cut  $S$  of  $G$ ,  $S \cap V(yCx - [x, y]) \neq \emptyset$
- ▶ Let  $R^+$  denote the set of non-negative real numbers, and  $w: V(G) \implies R^+$ . For  $H \subset G$ , we write  $w(H) = \sum_{v \in H} w(v)$ . Define  $w(\emptyset) = 0$ .
- ▶ **Theorem:-** Let  $(G, xCy)$  be a strong circuit graph, and let  $w: V(G) \implies R^+$ . Then  $G$  contains an  $x$ - $y$  path  $P$  such that  $\sum_{v \in (P-y)} [w(v)]^{\log_3 2} \geq [w(G-y)]^{\log_3 2}$
- ▶ **Corollary:-** If  $(G, C)$  is a circuit graph and  $e \in E(C)$ . Then  $G$  contains a cycle  $T$  through  $e$  such that  $|E(T)| \geq |V(G)|^{\log_3 2}$
- ▶ **Corollary:-** Let  $G$  be a 3-connected planar graph, and let  $e \in E(G)$ . Then  $G$  contains a cycle  $C$  through  $e$  such that  $|E(C)| \geq |V(G)|^{\log_3 2}$

# Literature Review

- ◀ In 1963, Moon and Moser implicitly made the following conjecture by giving 3-connected planar graphs with  $circ(G) \leq 9|V(G)|^{\log_3 2}$  [1].
- ◀ In 2002, Chen and Yu proved that a 3-connected planar graph on  $n$  vertices contains a cycle of length at least  $\Omega(n^{\log_3 2})$  [2].
- ◀ In 1931 Whitney proved that every 4-connected triangulation of the plane contains a hamiltonian cycle.[3]
- ◀ In 1956 Tutte proved that every 4-connected planar graph contains a hamiltonian cycle.[4]



# Problem Statement

- ◀ How many edge-disjoint cycles of length  $\Omega(n^{\log_3 2})$  can be guaranteed in a 3-connected planar graph.

# References

- ◀ Simple paths on polyhedra.,J. W. Moon, L. Moser,1963[1]
- ◀ Long Cycles in 3-Connected Planar graph Guantao Chen, Xingxing Yu,2002[2]
- ◀ H. Whitney,A theorem on graphs,Ann of math[3]
- ◀ T. Tutte,A Theorem on planar graph[4]