A Survey on Byzantine Gathering of Mobile Agents

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Outline



- 2 The Problem
- 3 Literature Survey



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- Gathering means meeting of two or more mobile agents starting from different positions in some topology.
- Termed as Rendezvous in case only two mobile agents are there.

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- Gathering means meeting of two or more mobile agents starting from different positions in some topology.
- Termed as Rendezvous in case only two mobile agents are there.
- Motivation:
 - How quickly two friends can meet in an unknown city.
 - Software agents collecting data from a computer network.

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Introduction | Inconsistency and Faults

• Inconsistency and faults are unavoidable in real life. Distributed computing for mobile agents on networks are no different.

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Introduction | Inconsistency and Faults

- Inconsistency and faults are unavoidable in real life. Distributed computing for mobile agents on networks are no different.
- Usually, two types of faults are considered:
 - Crash Faults: Similar to hardware failure[8].
 - **Byzantine faults:** May not obey the algorithm. Controlled by some adversary [4].
 - Weak Byzantine
 - Strong Byzantine

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f-Weak Byzantine Gathering: k mobile agents are dispersed on an anonymous port-labelled undirected graph G having n nodes. At most f agents among those k agents are weak Byzantine. The agents know n but do not know f, k or the ID range. The adversary wakes up at least one good agent. All the good agents need to gather at some node of the graph and terminate.

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Literature Survey

	Input	Byzantine	Condition of # Byzantine agents	Startup delay	Simultaneous termination	Time complexity
[Dieudonné et al.(2014)]	n	Weak	$f+1 \leq k$	Possible	Possible	$O(n^4 \cdot L \cdot X(n))$
[Dieudonné et al.(2014)]	f	Weak	$2f+2 \leq k$	Possible	Possible	Poly. of n and $ L $
[Bouchard et al.(2016)]	n, f	Strong	$2f+1 \leq k$	Possible	Possible	Exp. of n and $ L $
[Bouchard et al.(2016)]	f	Strong	$2f+2 \leq k$	Possible	Possible	Exp. of n and $ L $
[Hirose et al.(2021)]	N	Weak	$4f^2+9f+4\leq k$	Possible	No guarantee	$O((f+ L)\cdot X(N))$
[Hirose et al.(2021)]	N	Weak	$4f^2+9f+4\leq k$	Possible	Possible	$O((f+ \overline{L})\cdot X(N))$
[Hirose et al.(2022)]	N	Weak	$9f+8 \leq k$	Impossible	No guarantee	$O(f \cdot L \cdot X(N))$
[Saxena and Mondal(2024)]	n	Weak	$f^2+5f+9\leq k$	Possible	Possible	$O(k^2 \cdot L \cdot X(n))$

Table 1: An overview of synchronous Byzantine gathering algorithms with unique IDs. Here n is the number of nodes, N is the upper bound of n, l is the smallest ID among non-Byzantine agents, |L| is the length of the largest ID among non-Byzantine agents, $|\overline{L}|$ is the length of the largest ID among agents, k is the number of agents, X(n) is the number of rounds required to explore any network composed of n nodes, and f is the number of Byzantine agents.

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